

Static Analysis and code review

A journey through time

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Audit this code

```
/* Make these globals rather than local to mapping_chdir to avoid stack overflow */
char pathspace[MAXPATHLEN];
char old_mapped_path[MAXPATHLEN];
```

```
void do_elem(char *dir)
{
    /* . */
    if (dir[0] == '.' && dir[1] == '\0') {
        /* ignore it */
        return;
    }

    /* .. */
    if (dir[0] == '.' && dir[1] == '.' && dir[2] == '\0') {
        char *last;
        /* lop the last directory off the path */
        if ((last = strrchr(mapped_path, '/')) {
            /* If start of pathname leave the / */
            if (last == mapped_path)
                last++;
            *last = '\0';
        }
        return;
    }

    /* append the dir part with a leading / unless at root */
    if (!mapped_path[0] == '/' && mapped_path[1] == '\0')
        if (strlen(mapped_path) < sizeof(mapped_path) - 1)
            strcat(mapped_path, "/");
    if (sizeof(mapped_path) - strlen(mapped_path) > 1)
        strncat(mapped_path, dir, sizeof(mapped_path) - strlen(mapped_path) - 1);
}

int mapping_chdir(char *orig_path)
{
    int ret;
    char *sl, *path;

    strcpy(old_mapped_path, mapped_path);
    path = &pathspace[0];
    strcpy(path, orig_path);

    /* / at start of path, set the start of the mapped_path to */
    if (path[0] == '/') {
        mapped_path[0] = '/';
        mapped_path[1] = '\0';
        path++;
    }
}
```

```
/* Helper function for sgetpwnam(). */
char *sgetsave(char *s)
{
    char *new;

    new = (char *) malloc(strlen(s) + 1);

    if (new == NULL) {
        perror_reply(421, "Local resource failure: malloc");
        dologout(1);
        /* NOTREACHED */
    }
    (void) strcpy(new, s);
    return (new);
}

/* Save the result of a getpwnam. Used for USER command, since the data
 * returned must not be clobbered by any other command (e.g., globbing). */
struct passwd *sgetpwnam(char *name)
{
    static struct passwd save;
    register struct passwd *p;
#ifdef M_UNIX
    struct passwd *ret = (struct passwd *) NULL;
#endif
    char *sgetsave(char *s);
#ifdef KERBEROS
    register struct authorization *q;
#endif /* KERBEROS */

#ifdef SecureWare || defined(HPUX_10_TRUSTED)
    struct pr_passwd *pr;
#endif

#ifdef KERBEROS
    init_krb();
    q = getauthuid(p->pw_uid);
    end_krb();
#endif /* KERBEROS */

#ifdef M_UNIX
    if defined(SecureWare) || defined(HPUX_10_TRUSTED)
        if ((pr = getprpwnam(name)) == NULL)
            goto DONE;
#endif /* SecureWare || HPUX_10_TRUSTED */
    if ((p = getpwnam(name)) == NULL)
        goto DONE;
    #else /* M_UNIX */
    if defined(SecureWare) || defined(HPUX_10_TRUSTED)
        if ((pr = getprpwnam(name)) == NULL)
            return ((struct passwd *) pr);
    #endif

    #endif
```

A quick recap: Where we are today

- Small coding errors can have a big effect on security.
- Typical software development practices don't address the problem.
- As a group, developers tend to make the same security mistakes over and over.
- Static analysis tools can help identify common security errors early.

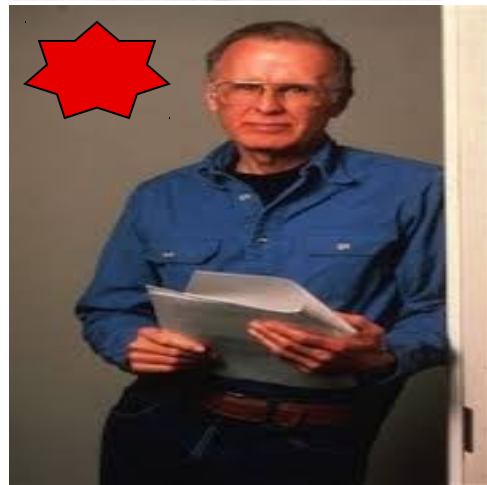


What to find?

- **Web-based enterprise applications**
 - ✓ **SQLI (and other injection attacks), XSS**
 - ✓ **Missing data validation**
 - ✓ **Session management errors (length, identifier)**
 - ✓ **Information leakage between sessions (concurrency)**
 - ✓ **Bad error handling (allowing system probing)**
 - ✓ **Compliance errors (treatment of personal data)**
 - ✓ **Insecure configuration**

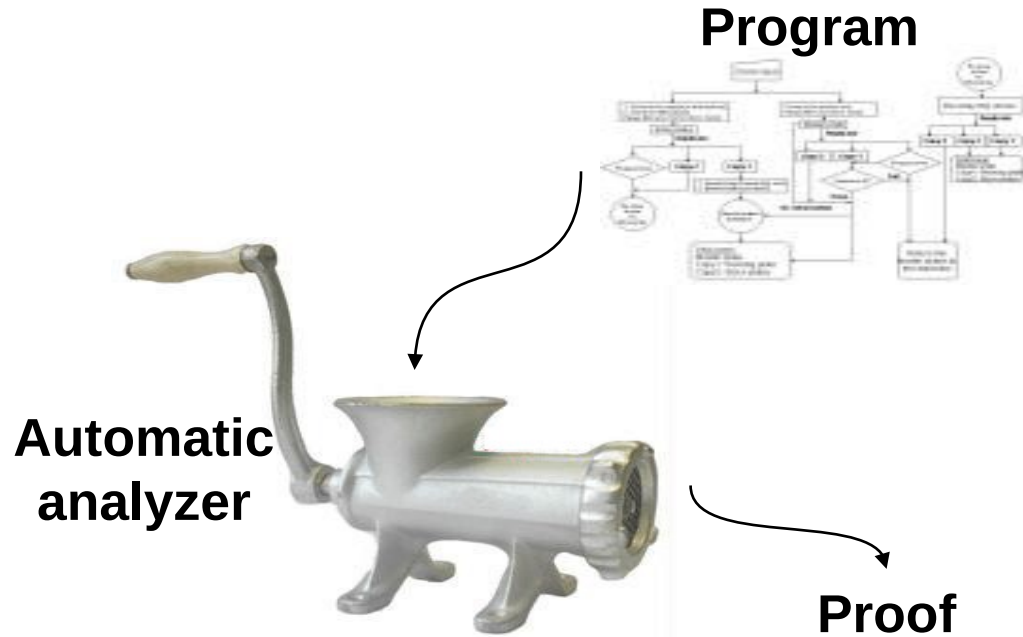






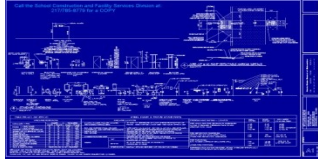
Proof

Program verification take #1



Program verification take #2

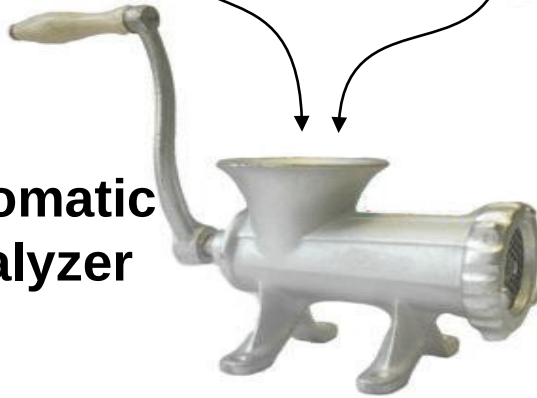
Specification



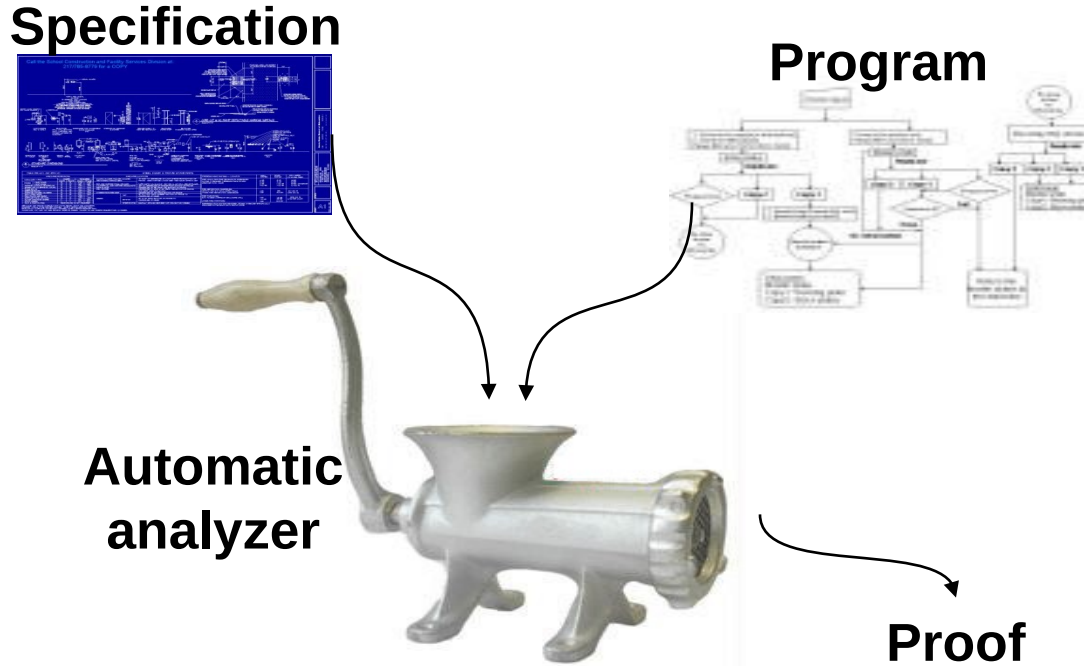
Program



**Automatic
analyzer**

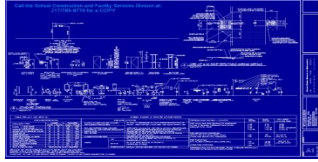


Proof



Program verification: a snag

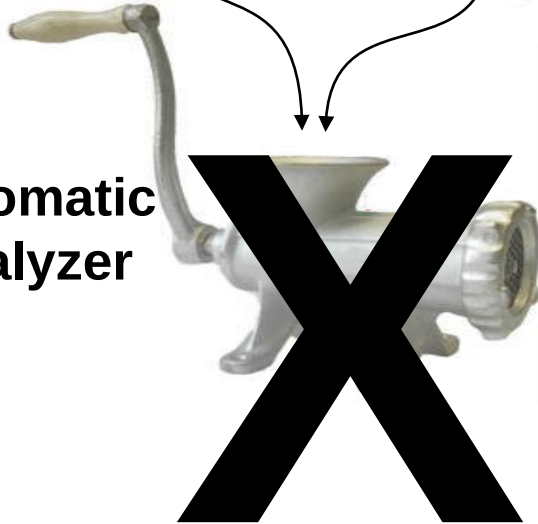
Specification



Program



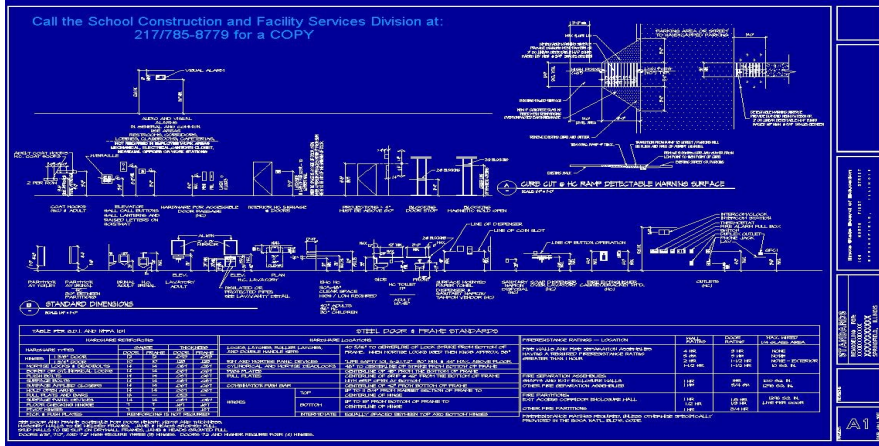
Automatic
analyzer



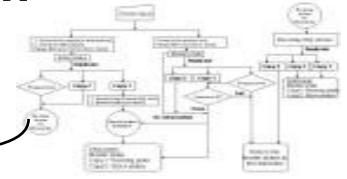
Proof

Program verification: a practical dilemma

Specification



Program

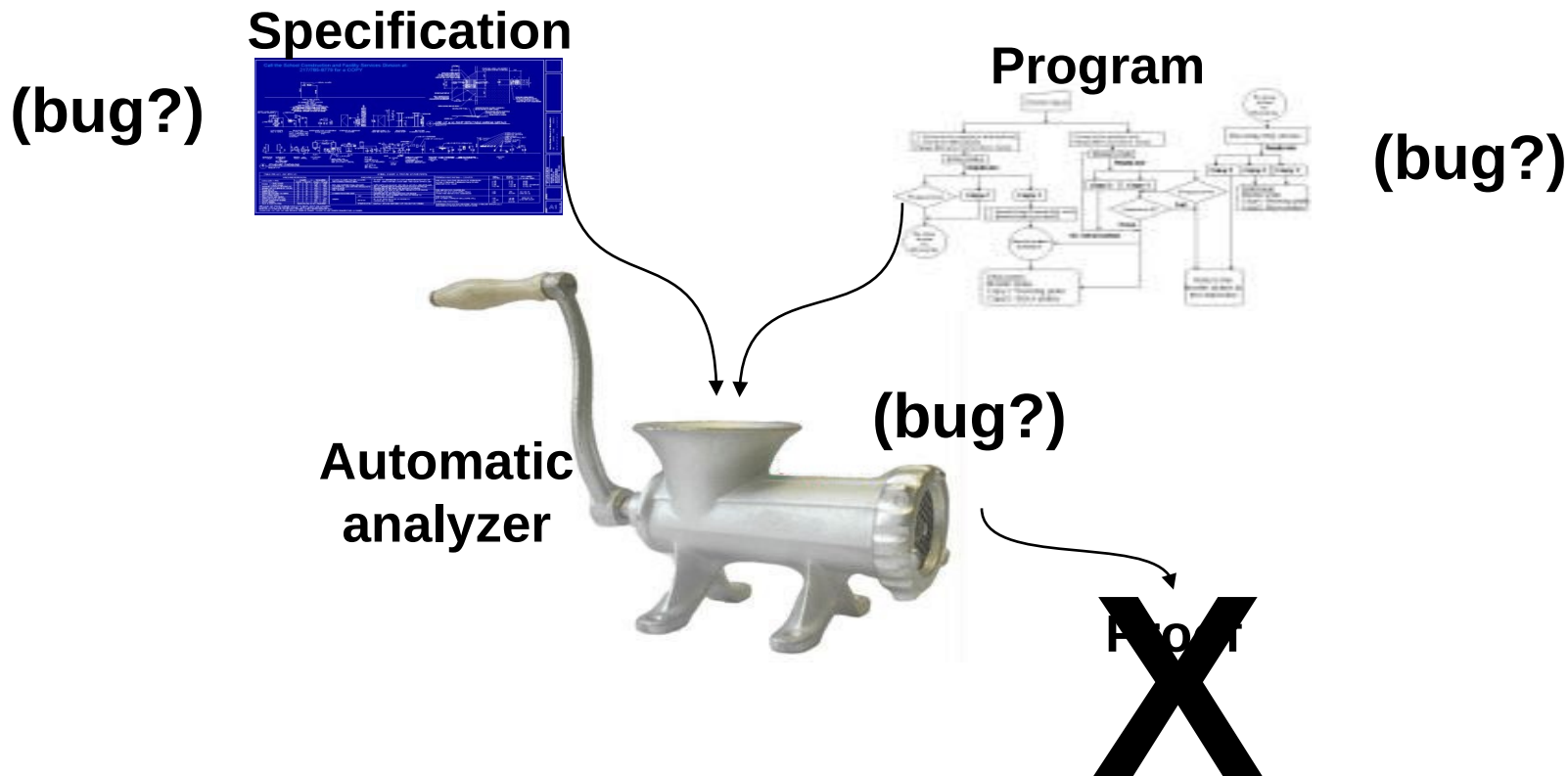


Automatic analyzer



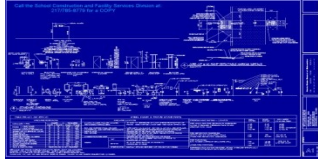
Proof

Program verification: the maiden voyage



Program verification: one more tweak

Specification



Program



**Automatic
analyzer**



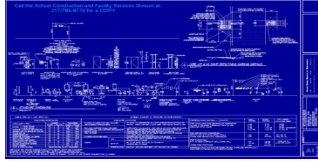
→
**Hints about what
went wrong
or
Proof**

time passes ...



Program verification today

Specification



Program



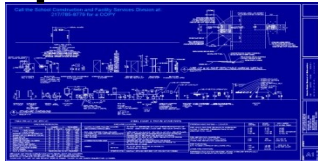
**Automatic
analyzer**



→
**Hints about what
went wrong
or
Proof**

A less ambitious plan

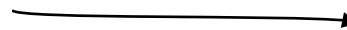
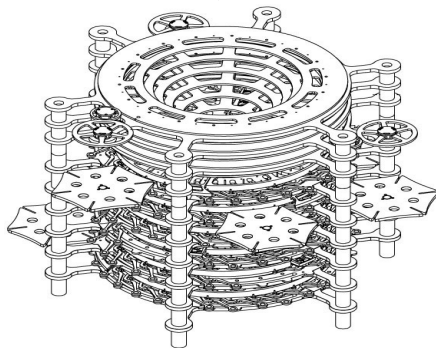
Specification



Program



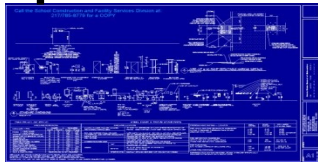
Checker



**Hints about what
went wrong
or
Proof**

A less ambitious plan

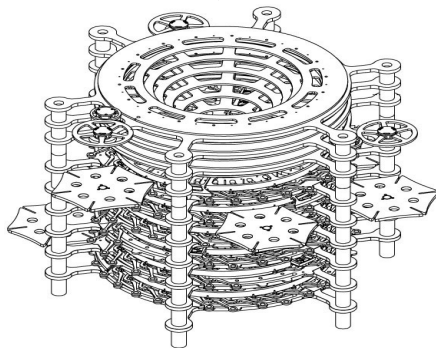
Specification



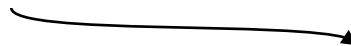
Program



Checker

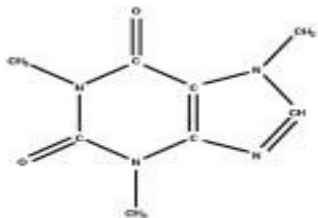


bug reports



A less ambitious plan

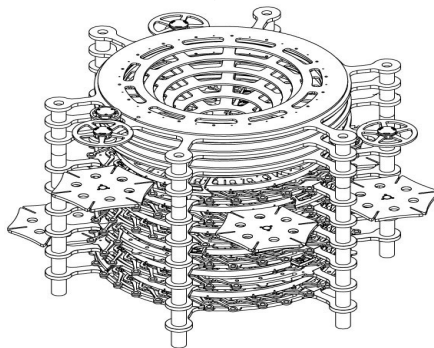
Property



Program

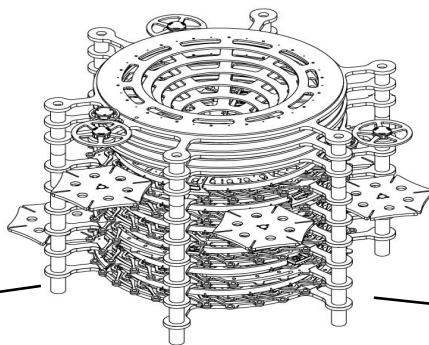


Checker



bug reports

Three hard problems



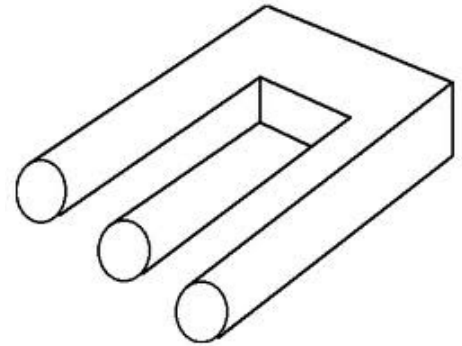
**Make
sense of
the
program**

**Compute
findings**

**Explain
findings to
user**

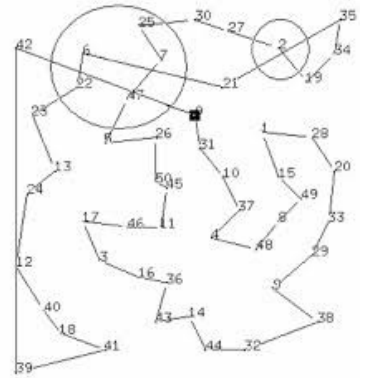
Make sense of the program

- Academic solutions typically target a limited number of languages (often tackling just a language subset).
- Enterprise applications use:
 - ✓ C/C++
 - ✓ Java / JSP / JS
 - ✓ PL/SQL
 - ✓ C#
 - ✓ Visual Basic
- Critical for success: robustness over precision



Compute findings

- Focus of most academic research
- Problem: No one-size fits all technique
- Solution: Build a flexible model, use multiple analyzers
- Gotchas: context, capacity
- Tricky:
 - ✓ pointers/pointer aliasing
 - ✓ function pointers/reflection/inversion of control (IOC)
 - ✓ loops



Context is King: token

read

Context is King: line

```
read(f, buf+len1, len-len1);
```


Context is King: function

```
static int my_read(int f, char *buf, int len)
{
    int len1 = 0;
    int ret;

    while(len1 < len) {
        ret = read(f, buf+len1, len-len1);
        if(ret < 0)
            return -1;
        len1 += ret;
    }
    return len;
}
```

Context is King: file

```
/* Make these globals rather than local to mapping_chdir to avoid stack overflow */
char pathspace[MAXPATHLEN];
char old_mapped_path[MAXPATHLEN];

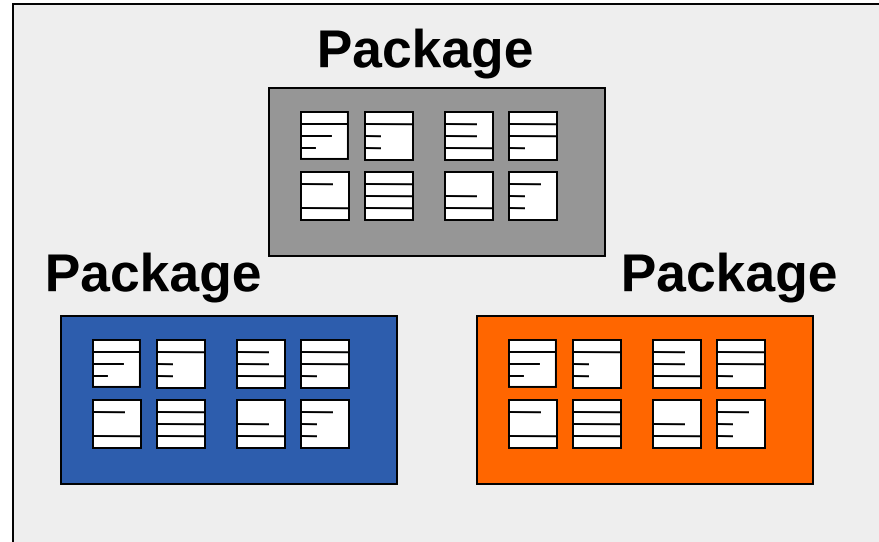
void do_elem(char *dir)
{
    /* . */
    if (dir[0] == '.' && dir[1] == '\0') {
        /* ignore it */
        return;
    }

    /* .. */
    if (dir[0] == '.' && dir[1] == '.' && dir[2] == '\0') {
        char *last;
        /* lop the last directory off the path */
        if ((last = strrchr(mapped_path, '/')) {
            /* If start of pathname leave the / */
            if (last == mapped_path)
                last++;
            *last = '\0';
        }
        return;
    }

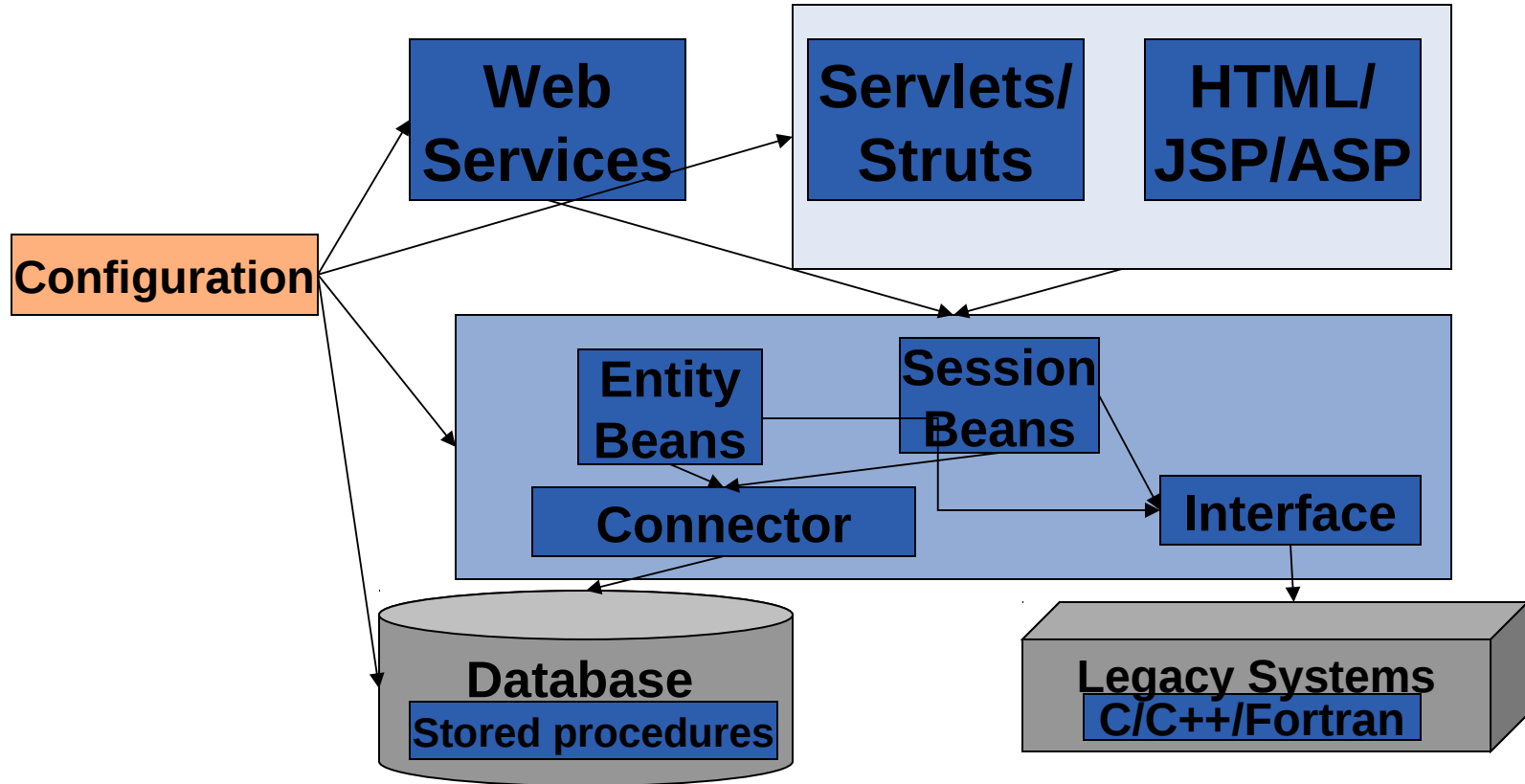
    /* append the dir part with a leading / unless at root */
    if (!(mapped_path[0] == '/' && mapped_path[1] == '\0'))
        if (strlen(mapped_path) < sizeof(mapped_path) - 1)
            strcat(mapped_path, "/");
    if (sizeof(mapped_path) - strlen(mapped_path) > 1)
        strncat(mapped_path, dir, sizeof(mapped_path) - strlen(mapped_path) - 1);
}
```

Context is King: process

Process



Context is King: system



Explain output to user

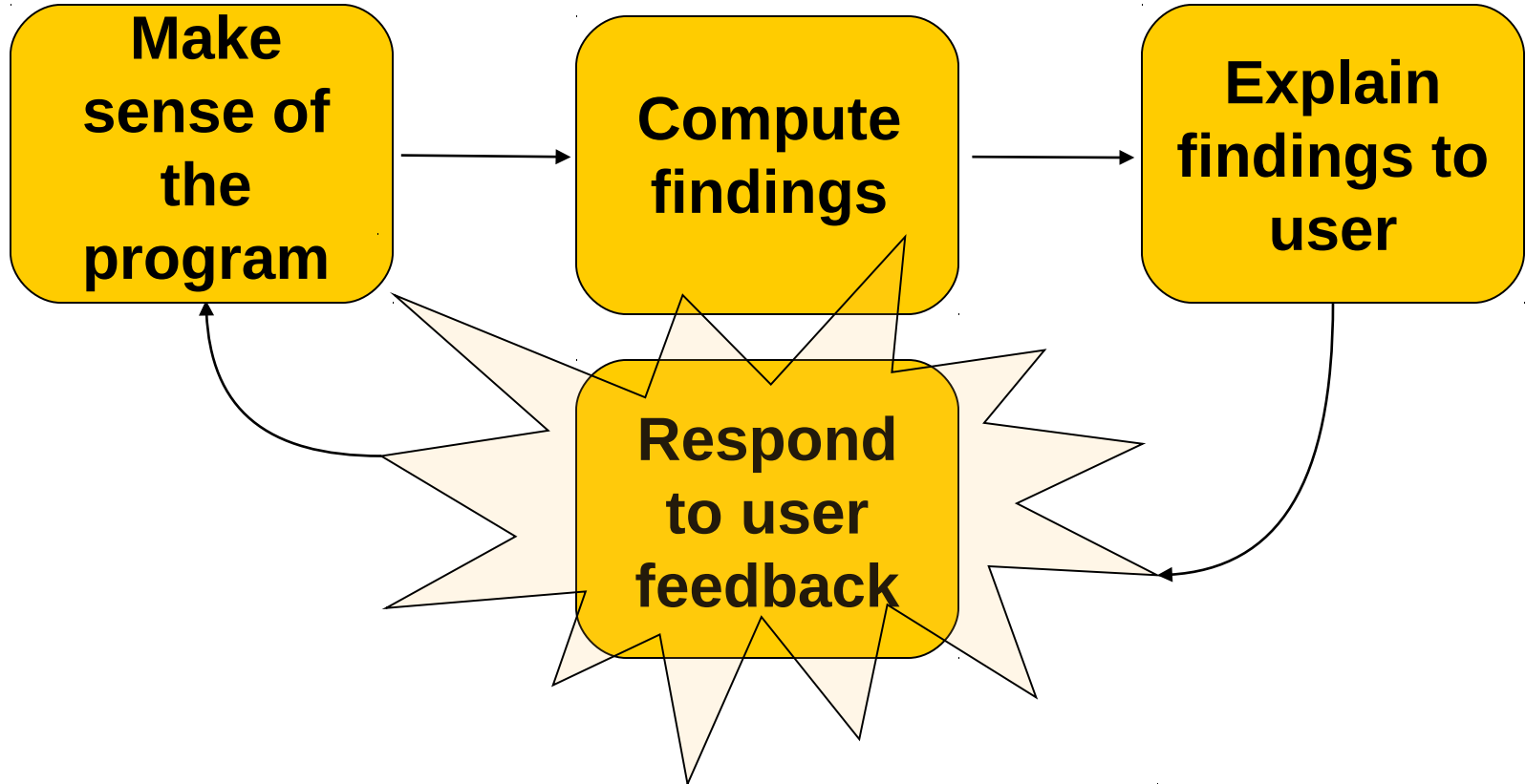
- Less studied in academia
- Most of the perceived value is actually here

The screenshot displays the Fortify Audit Workbench interface. At the top, there are tabs for 'Summary', 'Audit Guide', 'Scan', and 'Reports'. The main window is divided into several panes:

- Left Pane:** A 'Filter Set' section with 'Security Auditor View' selected. Below it, a 'Critical (673)' section is expanded to show a tree of issues. The selected issue is 'from StringUtil.java:49 (Privacy Violation)'. Below this is an 'Analysis Evidence' section showing a list of code snippets and their corresponding line numbers.
- Center Pane:** A code editor showing a snippet of JSP code. The line `<c:out value="${buttons}" escapeXml="false" />` is highlighted in orange, indicating the source of the violation.
- Right Pane:** A 'Functions' pane showing a tree of functions. The selected function is 'Assignment to userForm'.

Below the code editor, a 'Call Graph' is displayed. It shows the flow of control between different methods. The graph starts with 'StringUtil.encodePassword' (line 49) which calls 'UserAction.save'. This 'UserAction.save' method then calls 'Assignment to userForm' (line 177). The graph also shows other methods like 'setPassword', 'convert', 'copyProperties', and 'updateFormBean'.

A Critical Fourth Problem



Future developments

- **Use context to infer intent from the code**

Weakly typed languages

- **Speed up analysis utilizing multi core multi CPU setups**

Use the 'time gained' to do more analysis (breadth and depth)

- **Incremental analysis**

THANK YOU

GRACIAS
ARIGATO
SHUKURIA

THANKS
DANKSCHEEN
TASHAKKUR ATU
YASAWAYELAY
SUKSAMA
TRACI
BIYAN
SHUKRIA
GRAZIE
MEHRBANI
PALLERI
GODAMADITA
EFCAMBIETO
JUSPAKAR
BOLZIN
MERCİ