The Threat of Advanced Cross-Site Search Attacks







About me: Nethanel Gelernter

- Security Researcher / Hacker
 - Web application security
 - Ph.D., hacks, research papers, talks, etc.
- Cyberpion



- Exploring new attack vectors & developing defenses against some of them
- Leading the cyber-security studies & research in the College of Management, Israel



Agenda – practical timing attacks

- Cross-site search (XS-search) attacks & Response inflation
- Challenges
 - When response inflation is impossible
- Browser-based XS-search attacks
- Second-order XS-search attacks





Cross-Site Search Attacks

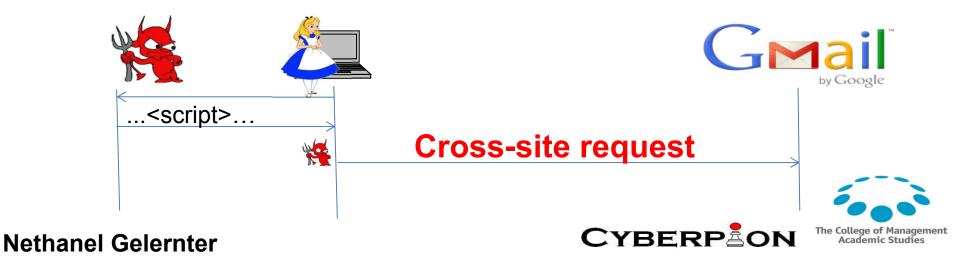
- Gelernter & Herzberg, CCS' 2015
- Exploit 'search' timing side-channel
- 'Search' in private-data kept by web-service
- Practical:
 - Tested on hundreds of Gmail users
 - Real world conditions
- Example: find user name
 - From lists of 2000 common (first and last) names
 - Takes about a minute





Cross-site attacker model

- Main model for web attacks
- The victim's browser is authenticated to services that hold private records (e.g., Gmail)
- The victim visits the attacker's website



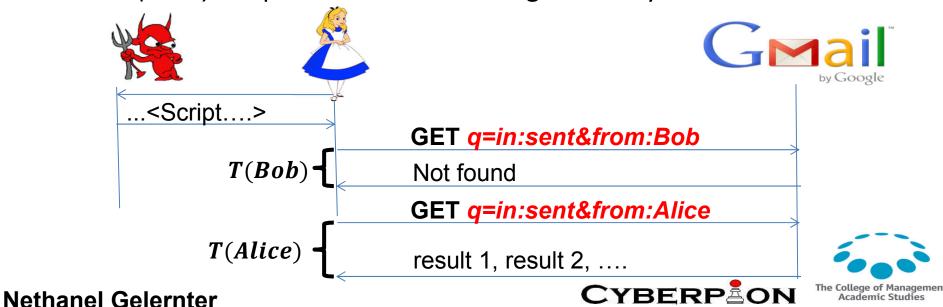
Cross-site attacker model

- Cross-site search over user's data in service
 - Attacker cannot access the content of the response
 - Same Origin Policy
 - The attacker can measure the response time (T)



XS-Search example: user name

- Find out whether the user is Alice or Bob...
- Compare:
 - T(Bob): response time for 'messages sent by Bob'
 - T(Alice): response time for 'messages sent by Alice'



What else can XS-Search expose?

Structured information









Email content

YAHOO!

XS-Search: Basic Flow

Find the answer for a Boolean question

- Three steps:
 - Transform the question into a search request
 - Send search requests and collect samples
 - Analyze response times

 answer the question!





XS-Search: Basic Flow – 1st Step

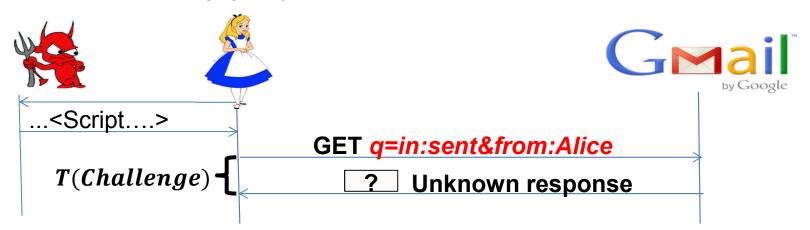
- Is the name of the user *Alice*?
 - in:sent from:Alice
- Is she related to bob@gmail.com?
 - bob@gmail.com&st=100
- Does Alice have an affair with Charlie
 - "I love you" to:Charlie from:Alice





XS-Search: Basic Flow – 2nd Step

- Send a **Challenge** request
 - Is the user name Alice?
 - True: a Full response is returned (has some content)
 - False: an **empty** response is returned

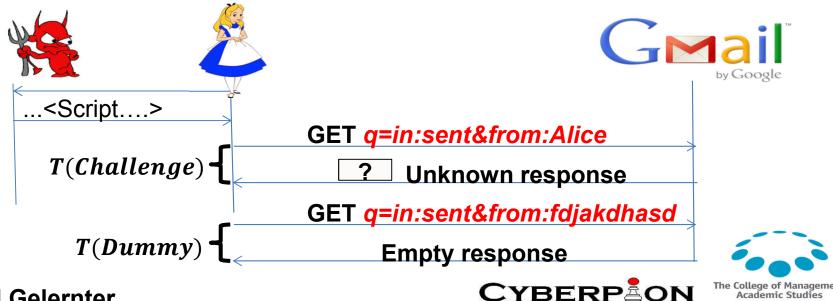




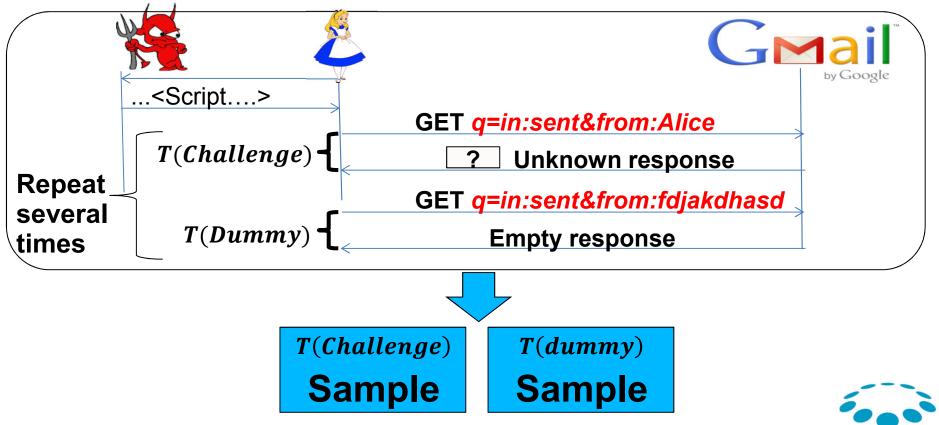


XS-Search: Basic Flow – 2nd Step

- Send a **Dummy** request
 - Is the user name fdjakdhasd?
 - The response is expected to be empty



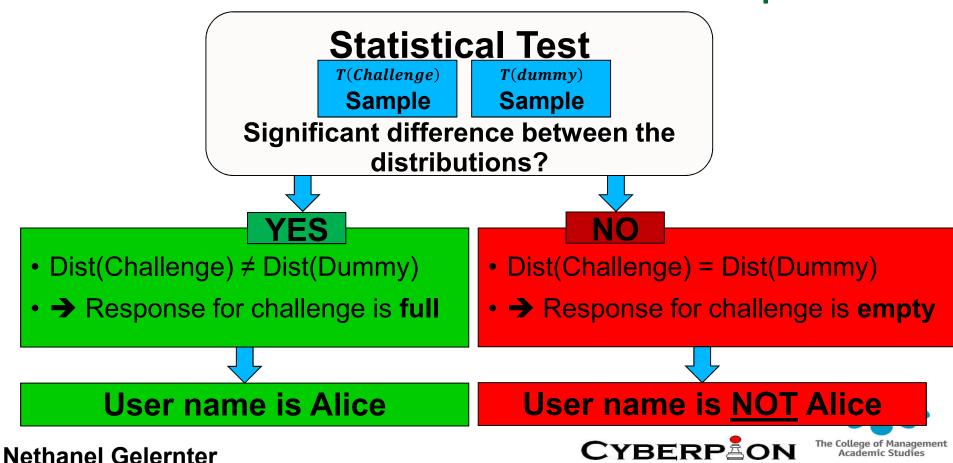
XS-Search: Basic Flow – 2nd Step







XS-Search: Basic Flow – 3rd Step



Practical timing attacks: challenges

- Timing attacks
 - Time measurements depend on dynamically-changing factors,
 e.g.: Congestion and concurrent processes in client and server
- Practical attacks
 - Minimal time
 - Exploit also short visits of users
 - Minimal number of requests
 - Avoid detection and blocking
 - E.g., by server's anti-DoS defenses





Response Inflation

- Increase the size difference between full and empty responses
- Larger difference in size → Larger difference in time

Larger → Slower





Response Inflation

- Search requests have many parameters
- Some of them are reflected in the responses as a function of the number of results

https://example.com/search?reflected_parameter=value

value

Empty response



Full response





Response Inflation

 Sometimes, the attacker can send very long strings as the value of the reflected parameter

https://example.com/search?reflected_parameter=Long string

Long string.....
Empty response

Response inflation example

- Exploiting Gmail search in the HTML view
- The query itself!
 - Appears once for each entry (50 max by default)
 - Can be inflated to 8KB
- Up to 400KB response size inflation!



Terms - Privacy - Gmail E

YBERPAON



But...







What if there is no response inflation?







What if there is no response inflation?

- Browser-based XS-search
 - When there is some difference in the response size
- Second-order XS-search
 - When there is **no** difference in the response size!





Browser-based (BB) XS-Search

- Statistical tests and divide and conquer algorithms
 - Gelernter & Herzberg, CCS' 2015
- Browser-based timing side channel
 - Van Goethem et al.,CCS' 2015
- Algorithmic improvements
 - Not in this talk







Classical vs. BB timing attacks

- Classical timing attacks:
 - Load the resources from the server several times to collect time measurements
- Browser-based timing attacks:
 - Load all the resources from the server once and cache them
 - Then load them from the cache many times to collect time measurements





Classical vs. BB timing attacks

- Exploiting / measurements affected by
 - Classical: network delay, server processing time, browser processing time
 - Browser-based: browser processing time
- Can be used to differentiate between
 - Classical: large/small resources, high/low server processing time
 - Browser-based: large/small resources





BB XS-Search: Basic Flow

Find the answer for a Boolean question

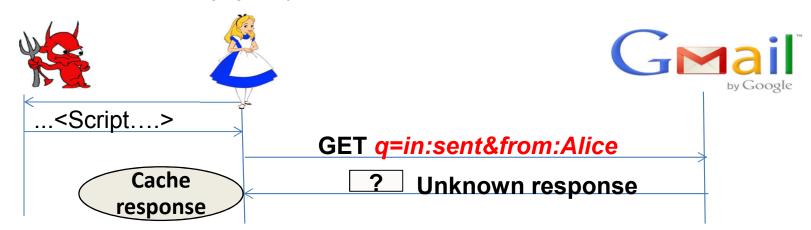
 Changing only the second step of the original XS-Search attack





BB XS-Search: Basic Flow – 2nd Step

- Send a **Challenge** request
 - Is the user name Alice?
 - True: a Full response is returned (has some content)
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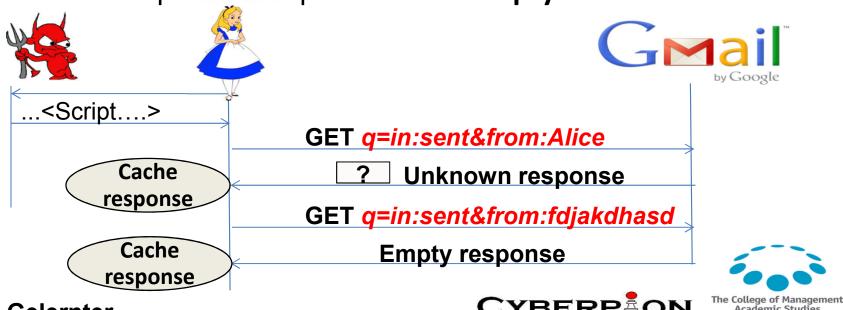






BB XS-Search: Basic Flow – 2nd Step

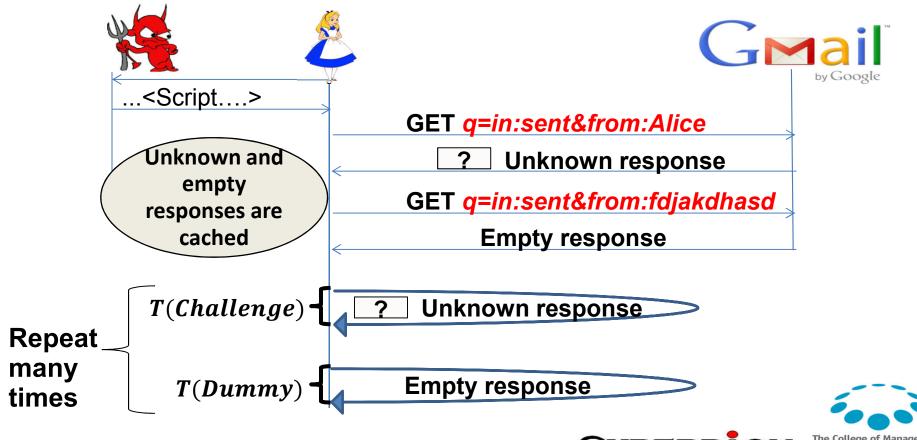
- Send a **Dummy** request
 - Is the user name fdjakdhasd?
 - The response is expected to be empty







BB XS-Search: Basic Flow – 2nd Step







Browser-based (BB) XS-Search

- Algorithmic improvements Not in this talk
- Not for Boolean questions
 - Basic flow only Boolean questions
 - Is the victim's name Alice?
- Answering multiple choice questions
 - E.g., which names out of many options are matching the victim?
- Optimally use the browser-based timing side-channel





Browser-based (BB) XS-Search

- Evaluation compared to both the previous works
- Repeating attacks/experiments done in each of them
 - Original XS-Search: extract victim's names from Gmail
 - BB timing attacks: extract victim's age from Facebook
- Significant improvement!
- In this talk: only one example





BB XS-Search vs. original XS-Search

- Gmail example
 - The goal of the attacker: extract the first and last names of the victim out of a list of 2000 names
 - XS-Search results:
 - 90% success rate (both first and last name found)
 - 1 minute on average
 - 2.6% false positive





BB XS-Search vs. original XS-Search

- Evaluation of the browser-based XS-search attack on 5 different Gmail accounts
 - 15-16 times on each of them
- Significant improvement!
 - 41.6 seconds on average (compared to 1 minute)
 - 92.3% success (compared to 89.7%)
 - 1.3% false positive (compared to 2.6%)



BB XS-Search vs. original XS-Search

DEMO





Second-order (SO) XS-Search attacks

- The problem: sometimes the size difference is negligible
- For example: a term that appears in a single email

Empty response

value

Full response





Second-order (SO) XS-Search attacks

- Second-order attacks
 - First, manipulate the attacked web application
 - Make it (more) vulnerable
 - Exploit the vulnerability
- Second-order XS-search attacks
 - First manipulate the attacked storage
 - Create significant response inflation
 - Launch browser-based XS-search attack





Second-order (SO) XS-Search attacks

- Two SO XS-search attacks
 - -Simple
 - —Inflating





Second-order (SO) XS-Search attacks

- Model
 - Storage
 - Many records
 - A secret appears in one of the records
- Attacker can manipulate the storage remotely
 - E.g., email accounts
 - Another example later...





Simple SO XS-Search attack

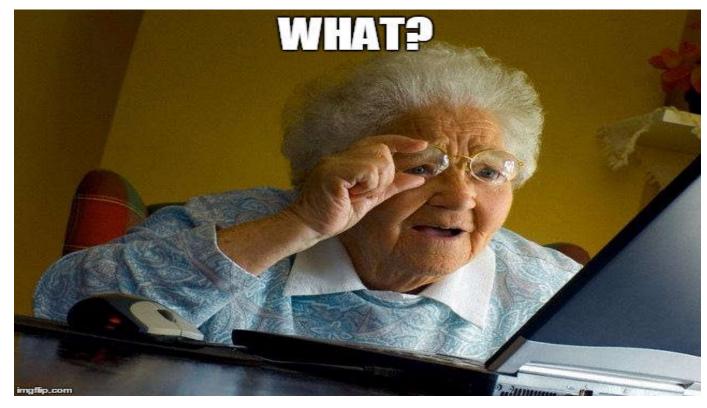
 The problem: the secret appears only once in the storage

 Simple solution: the attacker will add additional records that contain the secret!





Simple SO XS-Search attack



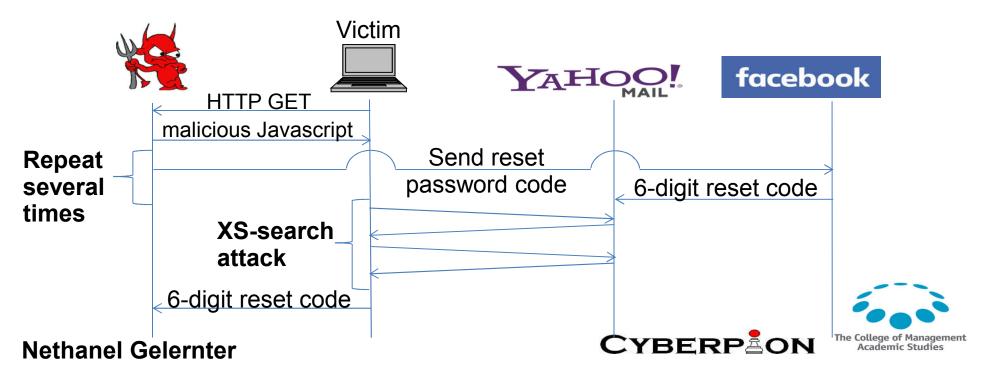






Simple SO XS-Search attack

 Example: extracting Facebook password-reset code from Yahoo! email



- Creates significant response inflation effect
 - Increases the size difference between empty and full response
- Unlike all the previous attacks: the empty response will be (significantly) larger than the full response





- The challenge of the attacker:
 - Find a secret out of a large dictionary of possible values
- Notations
 - M maximal number of results
 - Match-all record a record that contains all the possible values for the secret
 - Inflating record a record that significantly inflates the size of every response containing it



Attack process

First part:

- Plant one match-all inflating record in the storage
- Plant additional M-1 match-all records
- Additional record(s) may be added as a result of the victim's operations, or via operations triggered by the attacker

Second part:

Launch (browser-based) XS-search attack!







Response for searching the **right** secret New record (contains secret) Match-all record Match-all record M-1 Match-all record Inflating match-all record Response for searching the **wrong** secret





- Inflating record in email service providers
 - Email headers
 - From
 - To





- Example: extracting Visa/Mastercard credit card number
 - Structured information
 - VVVV-XXXX-YYYY-ZZZZ
- First and last names: extract 2 out of 2000
 - Done successfully!
- Credit card number: extract 4 out of 10000
 - Should not be much harder





- Example: extracting Visa/Mastercard credit card number
- Match-all record a record that contains all the possible 4-digit sequences
 - Possibly as an attachment
- Inflating match-all record a match-all record with very long From field



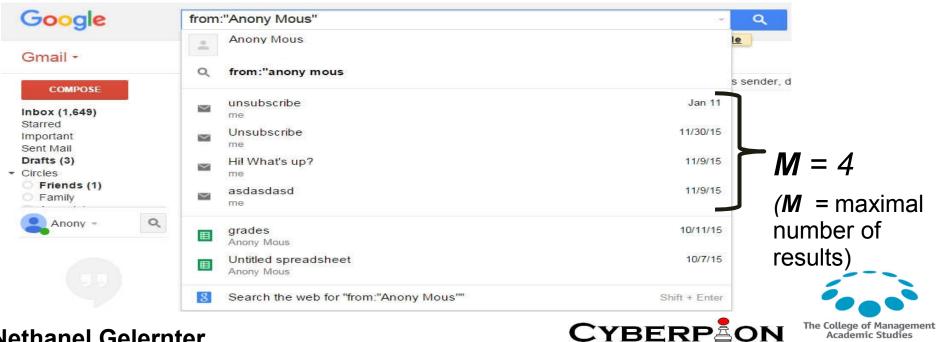


- Gmail example
- How?
 - Cross-site search requests are now blocked in both the HTML and standard views
- Cross-site search attack without sending cross-site search requests?



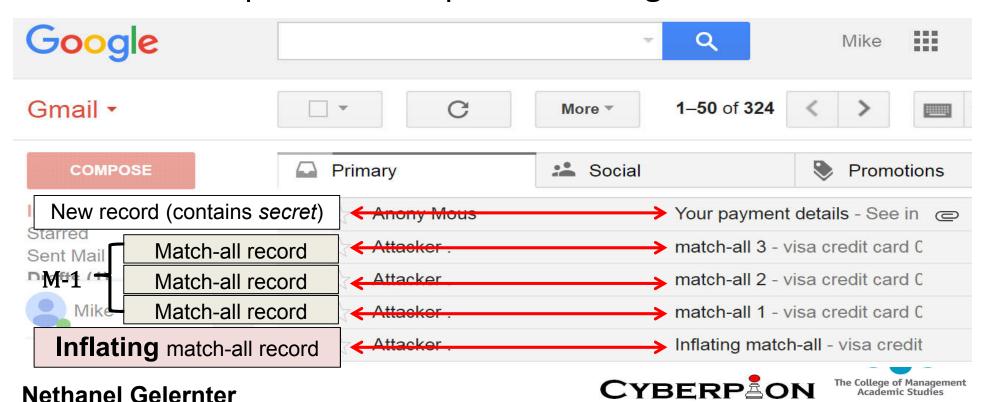


- Gmail example
- Exploiting the autocomplete feature!

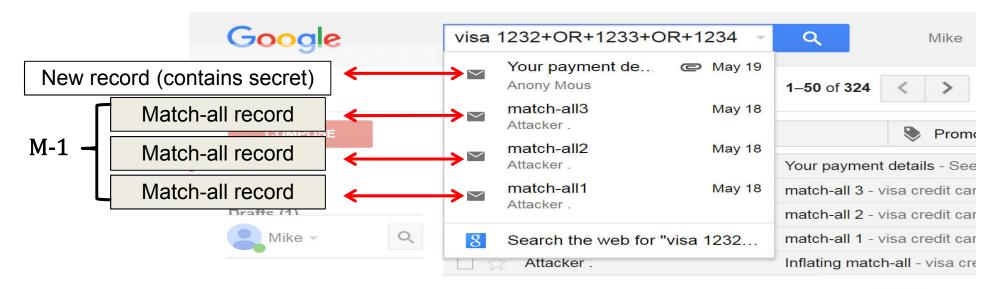


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Gmail example: the manipulated storage



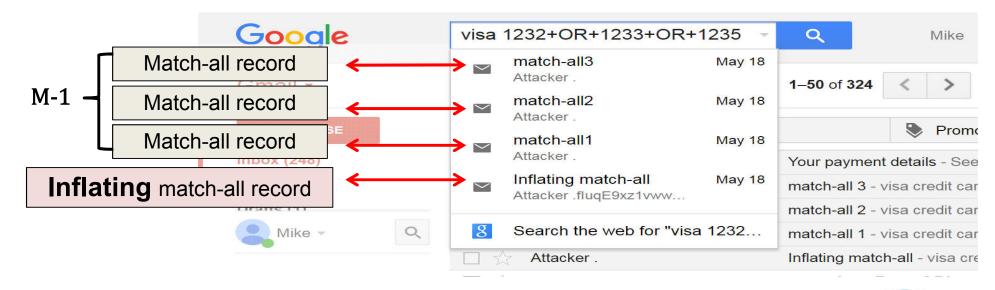
• Gmail example: full response (size is small)







• Gmail example: empty response (size is very large)







DEMO





- Evaluation results
 - 96% success rate within less than 50 seconds
 - Yet, in the other 4% percent, 3 out of 4 sequences were found, and it was possible to detect the error and to fix it





- The challenge: manipulations on the storage can be detected!
- Solution: manipulate the storage in a way that will not be detected by the user
- HOW?





- Emails solution: abuse anti-spam mechanisms
- The planted emails will be marked as spam
 - Users do not get notifications for spam emails
 - Users (usually) do not visit their spam folder
- Only when it is possible to search in the spam and in the other folders using the same request
 - E.g., Gmail
 - in:inbox OR in:spam

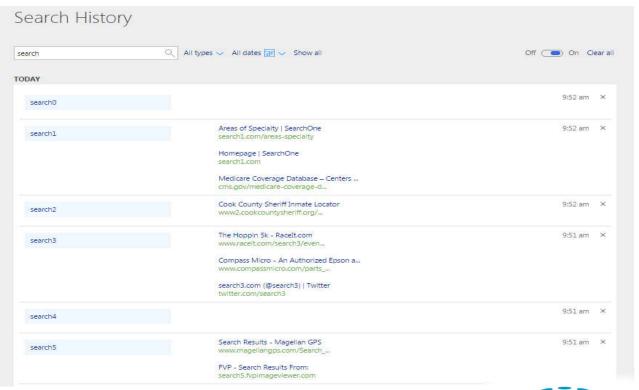


- Search history
- Two requirement for inflating SO XS-Search attack:
 - Inject records to the search history log
 - DONE: Gelernter & Grinstein & Herzberg, ACSAC 2015
 - Inject an inflating record





 Bing example: inflating SO XS-Search attack to extract search history







Defenses (briefly)

- If possible blocking cross-site search requests
- In other cases make it harder to exploit
 - Block inflation techniques
 - Rate limit
- Like (almost) every other web-application attack the challenge is to find all the vulnerable spots



Conclusions

- Advanced cross-site search attacks
 - Browser-based
 - Second order
- Practical!
- Many vulnerable websites
 - Including popular ones





Thank you!





Questions?



